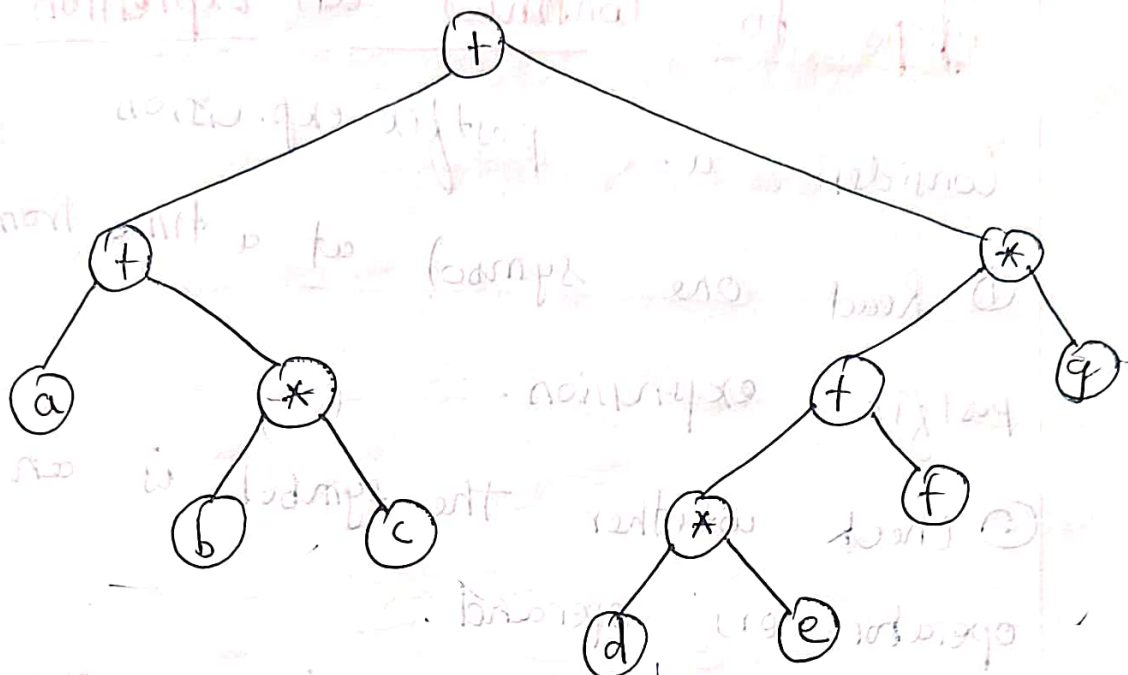


## Expression Tree

The expression tree for the expression  $(a + b * c) + ((d * e + f) * g)$  will be shown in following figure



\* The leaves of an expression tree are operands, such as constants or variable names, and the other nodes contain operators

\* The tree happens to be binary; because all the operations are binary. we can evaluate an expression tree  $T$ , by applying the operator at the root to the values obtained by recursively evaluating the left and right subtrees.

\* In above example the left subtree evaluates  $a + (b * c)$  and right subtree evaluates  $((d * e + f) * g)$

## Steps to Construct an Expression Tree

Consider a postfix expression

① Read one symbol at a time from postfix expression.

② Check whether the symbol is an operator or operand.

(a) If the symbol is an operand create one node and push a pointer on to the stack.

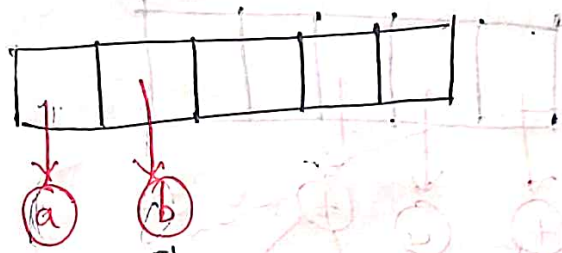
(b) If the symbol is an operator pop two pointers from the stack namely  $T_1$  &  $T_2$  and form a new tree with operator as the root and  $T_1$  as left child &  $T_2$  as right child. pointer to new tree is pushed on to the stack.

### Example

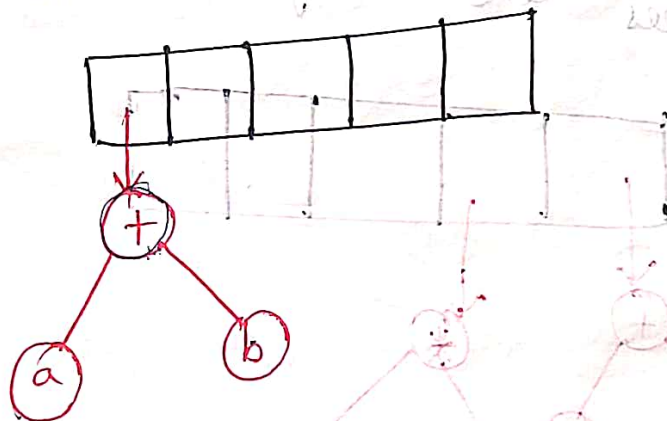
a b + c d e + \* \*



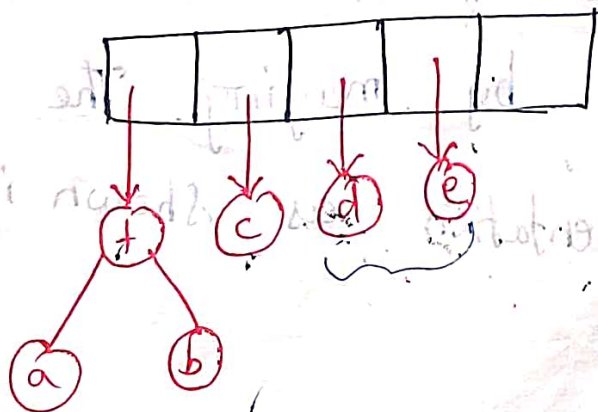
① First two symbols are operands, so we can create one node tree and push their pointers on to the stack.



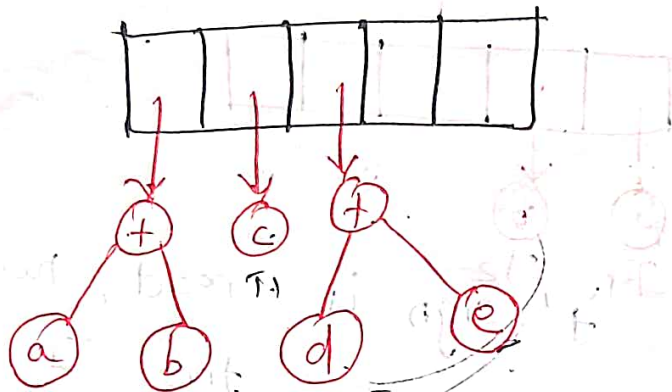
② Next '+' sign is read, two pointers are popped and a new tree is formed and a pointer is again pushed on to the stack.



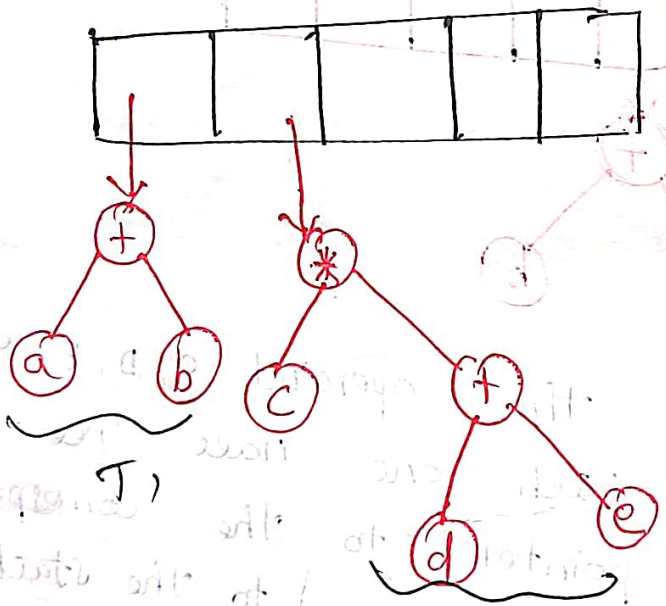
③ Next the operand c, d, e are read and for each one node tree is created and a pointer to the corresponding tree is pushed on to the stack.



(4) Now the '+' is read so last two trees are merged.



(5) Now the '\*' is read, so the last two trees are merged.



(6) Now the 'x' is read and tree is formed by merging. The final representation as shown in figure.

