

## Conversion from NFA to DFA

In this section, we will discuss the method of converting NFA to its equivalent DFA. In NFA, when a specific input is given to the current state, the machine goes to multiple states. It can have zero, one or more than one move on a given input symbol. On the other hand, in DFA, when a specific input is given to the current state, the machine goes to only one state. DFA has only one move on a given input symbol.

Let,  $M = (Q, \Sigma, \delta, q_0, F)$  is an NFA which accepts the language  $L(M)$ . There should be equivalent DFA denoted by  $M' = (Q', \Sigma', q_0', \delta', F')$  such that  $L(M) = L(M')$ .

Steps for converting NFA to DFA:

**Step 1:** Initially  $Q' = \phi$

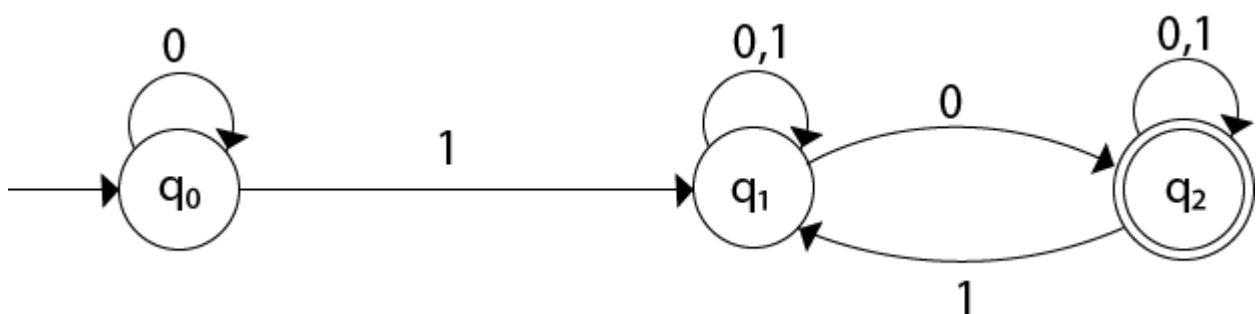
**Step 2:** Add  $q_0$  of NFA to  $Q'$ . Then find the transitions from this start state.

states is not in  $Q'$ , then add it to  $Q'$ .

**Step 4:** In DFA, the final state will be all the states which contain  $F$  (final states of NFA)

Example 1:

Convert the given NFA to DFA.



**Solution:** For the given transition diagram we will first construct the transition table.

| State             | 0              | 1              |
|-------------------|----------------|----------------|
| $\rightarrow q_0$ | $q_0$          | $q_1$          |
| $q_1$             | $\{q_1, q_2\}$ | $q_1$          |
| $*q_2$            | $q_2$          | $\{q_1, q_2\}$ |

Now we will obtain  $\delta'$  transition for state  $q_0$ .

1.  $\delta'([q_0], 0) = [q_0]$
2.  $\delta'([q_0], 1) = [q_1]$

The  $\delta'$  transition for state  $q_1$  is obtained as:

#### ADVERTISEMENT

1.  $\delta'([q_1], 0) = [q_1, q_2]$  (new state generated)
2.  $\delta'([q_1], 1) = [q_1]$

The  $\delta'$  transition for state  $q_2$  is obtained as:

1.  $\delta'([q_2], 0) = [q_2]$
2.  $\delta'([q_2], 1) = [q_1, q_2]$

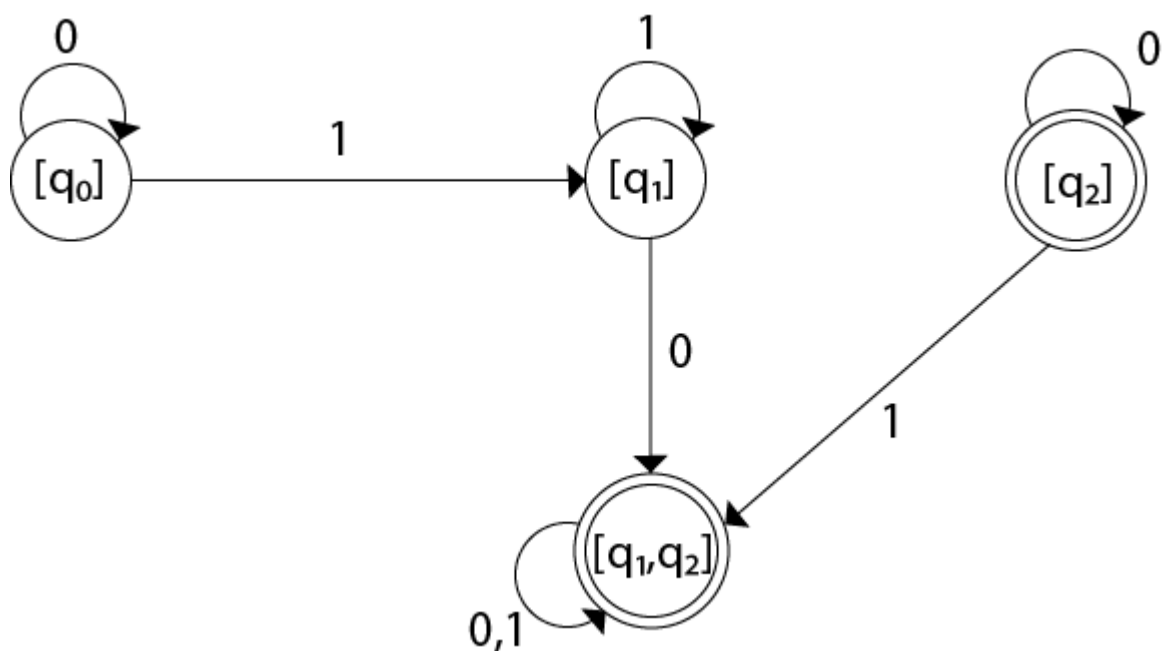
Now we will obtain  $\delta'$  transition on  $[q_1, q_2]$ .

1.  $\delta'([q_1, q_2], 0) = \delta(q_1, 0) \cup \delta(q_2, 0)$
2.  $\quad\quad\quad = \{q_1, q_2\} \cup \{q_2\}$
3.  $\quad\quad\quad = [q_1, q_2]$
4.  $\delta'([q_1, q_2], 1) = \delta(q_1, 1) \cup \delta(q_2, 1)$
5.  $\quad\quad\quad = \{q_1\} \cup \{q_1, q_2\}$
6.  $\quad\quad\quad = \{q_1, q_2\}$
7.  $\quad\quad\quad = [q_1, q_2]$

The state  $[q_1, q_2]$  is the final state as well because it contains a final state  $q_2$ . The transition table for the constructed DFA will be:

| State              | 0            | 1            |
|--------------------|--------------|--------------|
| $\rightarrow[q_0]$ | $[q_0]$      | $[q_1]$      |
| $[q_1]$            | $[q_1, q_2]$ | $[q_1]$      |
| $*[q_2]$           | $[q_2]$      | $[q_1, q_2]$ |
| $*[q_1, q_2]$      | $[q_1, q_2]$ | $[q_1, q_2]$ |

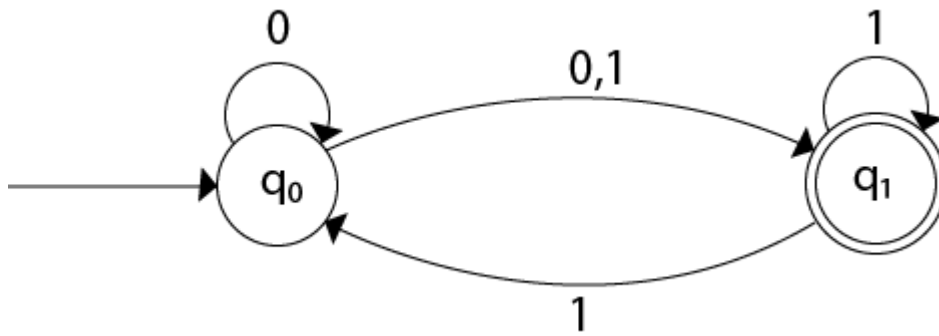
The Transition diagram will be:



The state  $q_2$  can be eliminated because  $q_2$  is an unreachable state.

Example 2:

Convert the given NFA to DFA.



**Solution:** For the given transition diagram we will first construct the transition table.

| State | 0        | 1        |
|-------|----------|----------|
| →q0   | {q0, q1} | {q1}     |
| *q1   | ϕ        | {q0, q1} |

Now we will obtain  $\delta'$  transition for state q0.

1.  $\delta'([q0], 0) = \{q0, q1\}$
2.  $\quad\quad\quad = [q0, q1]$  (**new state generated**)
3.  $\delta'([q0], 1) = \{q1\} = [q1]$

The  $\delta'$  transition for state q1 is obtained as:

1.  $\delta'([q1], 0) = \phi$
2.  $\delta'([q1], 1) = [q0, q1]$

Now we will obtain  $\delta'$  transition on  $[q0, q1]$ .

1.  $\delta'([q0, q1], 0) = \delta(q0, 0) \cup \delta(q1, 0)$
2.  $\quad\quad\quad = \{q0, q1\} \cup \phi$
3.  $\quad\quad\quad = \{q0, q1\}$
4.  $\quad\quad\quad = [q0, q1]$

Similarly,

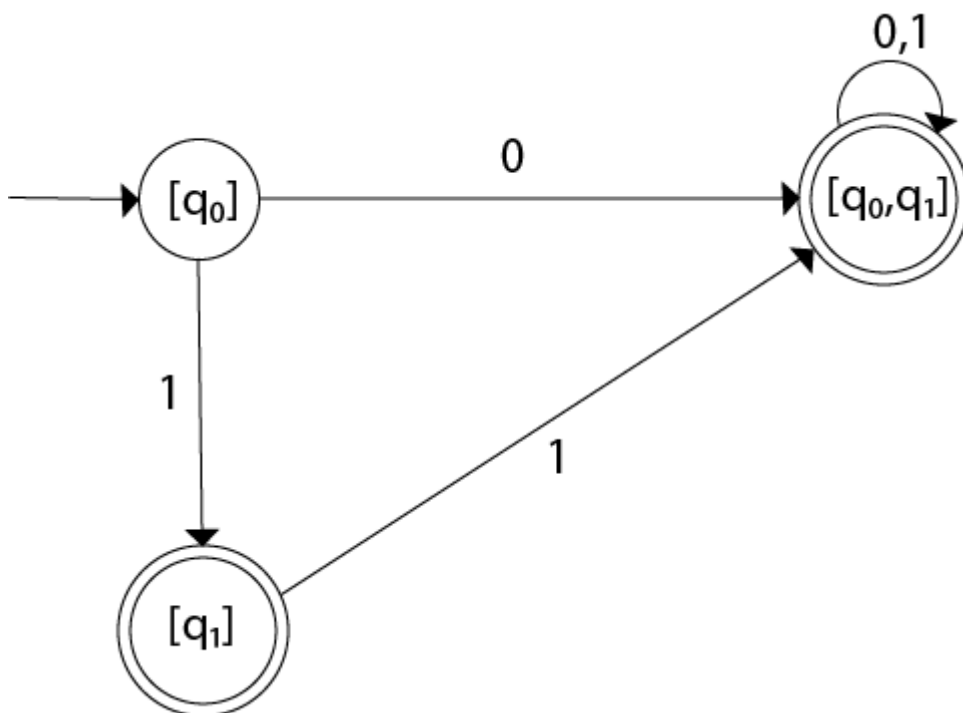
1.  $\delta'([q_0, q_1], 1) = \delta(q_0, 1) \cup \delta(q_1, 1)$
2.  $= \{q_1\} \cup \{q_0, q_1\}$
3.  $= \{q_0, q_1\}$
4.  $= [q_0, q_1]$

As in the given NFA,  $q_1$  is a final state, then in DFA wherever,  $q_1$  exists that state becomes a final state. Hence in the DFA, final states are  $[q_1]$  and  $[q_0, q_1]$ . Therefore set of final states  $F = \{[q_1], [q_0, q_1]\}$ .

The transition table for the constructed DFA will be:

| State              | 0            | 1            |
|--------------------|--------------|--------------|
| $\rightarrow[q_0]$ | $[q_0, q_1]$ | $[q_1]$      |
| $*[q_1]$           | $\phi$       | $[q_0, q_1]$ |
| $*[q_0, q_1]$      | $[q_0, q_1]$ | $[q_0, q_1]$ |

The Transition diagram will be:



Even we can change the name of the states of DFA.

**Suppose**

1. A = [q0]
2. B = [q1]
3. C = [q0, q1]

With these new names the DFA will be as follows:

