

SNS COLLEGE OF ENGINEERING Coimbatore - 641 107



TOPIC:11.- PROOF METHODS AND STRATEGY.

The Theory of Inference for Predicate calculus Universal Specification (US) If a statement of the form (#x) [A(x) is assumed to be true, then the universal quantifier can be dropped to obtain A(y) is the for any arbitrary object 'y' in the universe. Existential Specification (ES) From Fr (A(2)) one can conclude A(y), provided that y is not free in any given premise and also not free in any prior step of the divivation Universal Generalization (UG) Friom A (2) one can / con lude $A(y) \Rightarrow (\forall x)(A(x))$ Existential Generalization (EG) $A(y) \Rightarrow (\exists x)(A(x))$

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Fxpress The Statement "Every student in this class
as completed Assignment - I' as a quantifiers
Let
$$C(x)$$
 : x is in this class
 $A(x)$: x has completed Assignment - I
For all x , if x is in this class, then x has
completed Assignment - 1.
For all x , if x is in this class, then x has
multid Assignment - 1.
Its symbolic form $(+x)(C(x) \rightarrow A(x))$

2.





Show that
$$(\forall x)(P(x) \rightarrow Q(x)) \land (\forall x)(Q(x) \rightarrow R)$$

$$\Rightarrow (\forall x)(P(x) \rightarrow R(x))$$

$1) (\forall x) (P(x) \rightarrow Q(x))$	Rule P
2) $P(y) \rightarrow Q(y)$	Rule US
3) $(\forall x)(Q(x) \rightarrow R(x))$	Rule P
	Rule US
	Rule T (P→Q,Q→R
$6)(\forall x)(P(x) \rightarrow R(x)).$	⇒ P→R) Rule UG
	2) $P(y) \rightarrow Q(y)$ 3) $(\forall x) (Q(x) \rightarrow R(x))$ 4) $Q(y) \rightarrow R(y)$ 5) $P(y) \rightarrow R(y)$



4)

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Show that $(\forall x) (P(x) \lor Q(x)) \Rightarrow (\forall x) P(x) \lor (\exists x)$ We use indirect method, by assuming $P(x) \lor (\exists x)(Q(x))$ as an additional promise

513	$ 1\rangle \neg \left[(\forall x) P(x) \lor (\exists x) Q(x) \right]$	Rule P
513	2) $(\exists z) \neg P(z) \land$ $(\forall z) \neg Q(z)$	Rule T (Demorgan's)
113	3) (3x) - p(x)	Rule T (PAQ ⇒ P)
513	4) (¥71) ¬Q(7)	Ruli T (PAQ ⇒Q)
513	5) - P(y)	Rule ES
513	6) - Q(4)	Rule US
513	7) - P(y) ~ - Q(y)	Rule T (P,Q ⇒ PAQ)
513	8) $\neg (P(y) \vee Q(y))$	Rule T (Demorgan's)
<u></u> {9}	9) $(\forall x) (P(x) \vee Q(x))$	Rule P

SNSCE/ S&H/ UNIT 1/ DM/11 – PROOF METHODS AND STRATEGY /MATHS

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<i>§1,9</i>	11) [P(y) VQ(y)] ^ - [P(y) VQ(y)]	Rule T (P, $a \Rightarrow PAQ$)	
\$1.97	12) F		

5) Verify the validity of the following argument. From living thing is a plant or animal John's gold fishing alive and it is not a plant. All animals have hearts. Thurefore John's gold fish has a heart. Let L(x): x is a living thing P(x): x is a plant A(x): x is an animal H(x) : x has a heart Then, the given premises are (1) $(\forall x) [L(x) \rightarrow (P(x) \cdot V A(x))]$ (2) L(j) A - P(j) (3) $(\forall x) [A(x) \rightarrow H(x)]$ conclusion is H(j)



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ALC: NO		
113	$) (\forall x) \left[L(x) \rightarrow (P(x) \lor A(x) \right]$	x)) Rule P
113	2) $L(j) \rightarrow P(j) \vee A(j)$) Rule US
{ 3 }	3) L(j) A¬P(j)	Rule P
{3 }	4) L(j), ¬P(j)	Rule T (PAQ ⇒P.Q)
{1,3 }	5) P(j) V {A(j)	Rule T(P, P→a ⇒ a)
51,3}	$6) \neg P(j) \rightarrow A(j)$	Rule T (P→a => ¬PVQ)
57}	$(\forall x) (A(x) \rightarrow H(x))$	Rule P
יי ר זי ר	8) $A(j) \rightarrow H(j)$	Rule US
ין 1'ז 1'ז	$q) \neg P(j) \rightarrow \mathbf{M}(j)$	Rule T (P→Q,Q→R ⇒ P→R)
1,3,7 {	10) H(j)	Rule T (P, P→a⇒a
10.17		