



# SNS COLLEGE OF ENGINEERING

Kurumbapalayam (Po), Coimbatore - 641 107

**AN AUTONOMOUS INSTITUTION**

Accredited by NAAC - UGC with 'A' Grade

Approved by AICTE, New Delhi & Affiliated to Anna University, Chennai

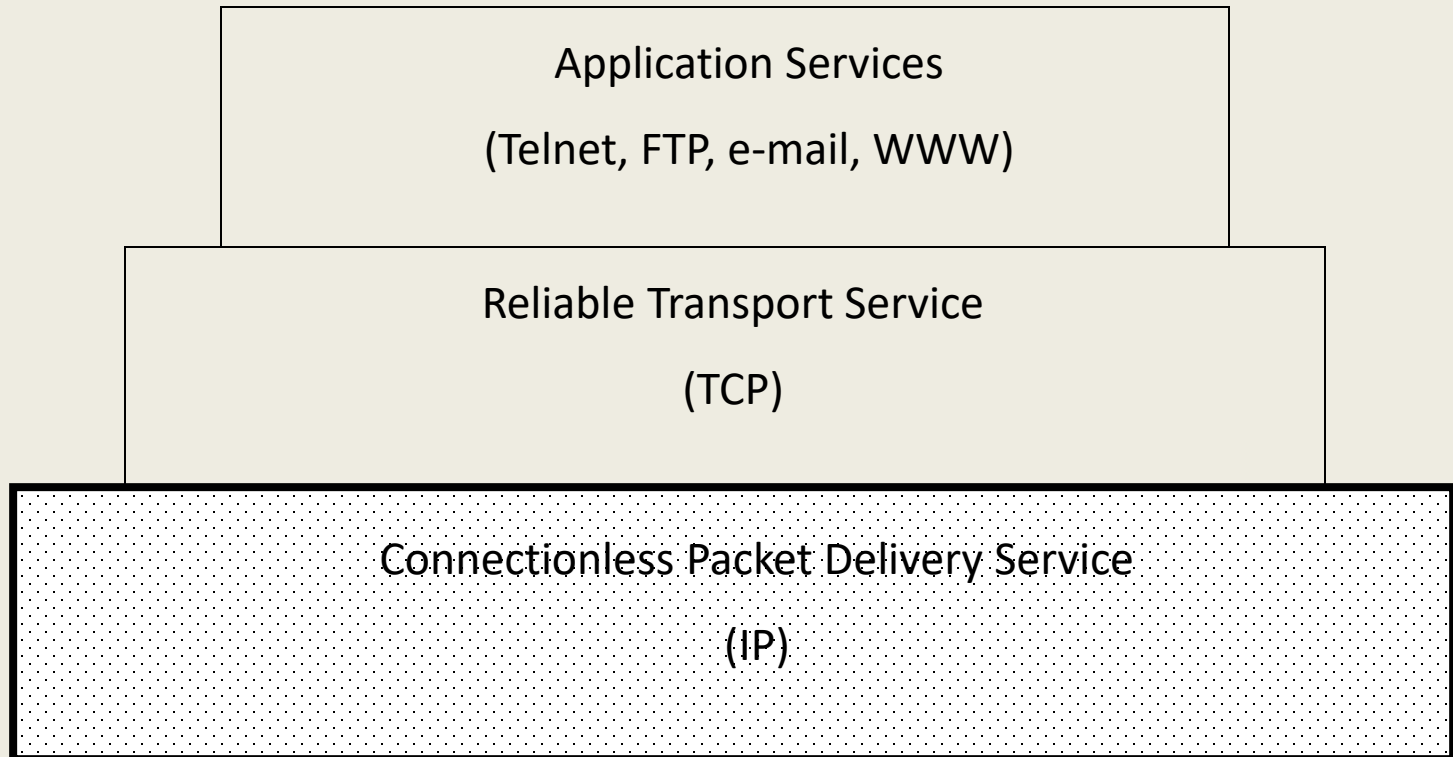


**MOBILE**  
CLOUD COMPUTING



# INTRODUCTION TO MOBILE COMPUTING

# The Internet Protocol (IP)

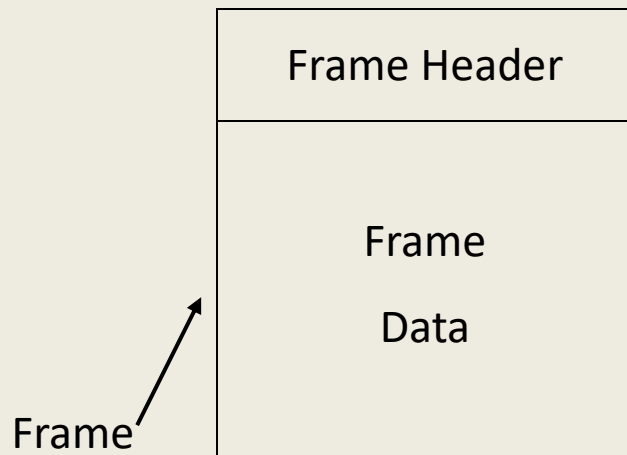


# The Internet Protocol (IP)

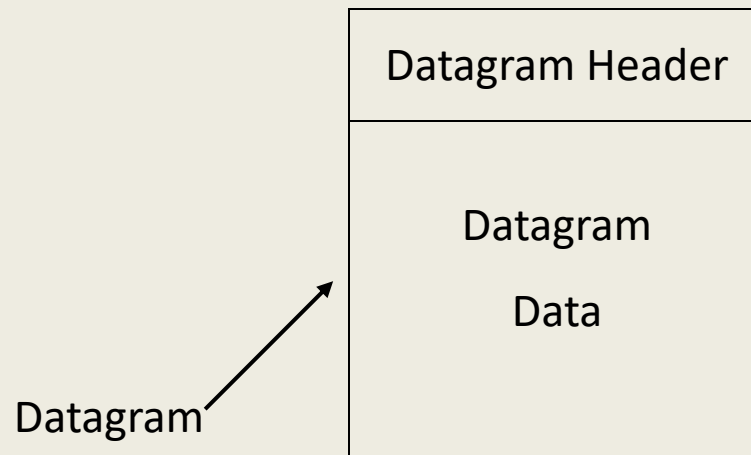
- Provides a packet delivery service which is:
  - Unreliable
  - Best-effort
  - Connectionless
- Defines the basic unit of data transfer
- Performs the routing function
- Includes a set of rules that embody the idea of unreliable packet delivery

# IP Datagrams

- The basic unit of data transfer



Physical Network

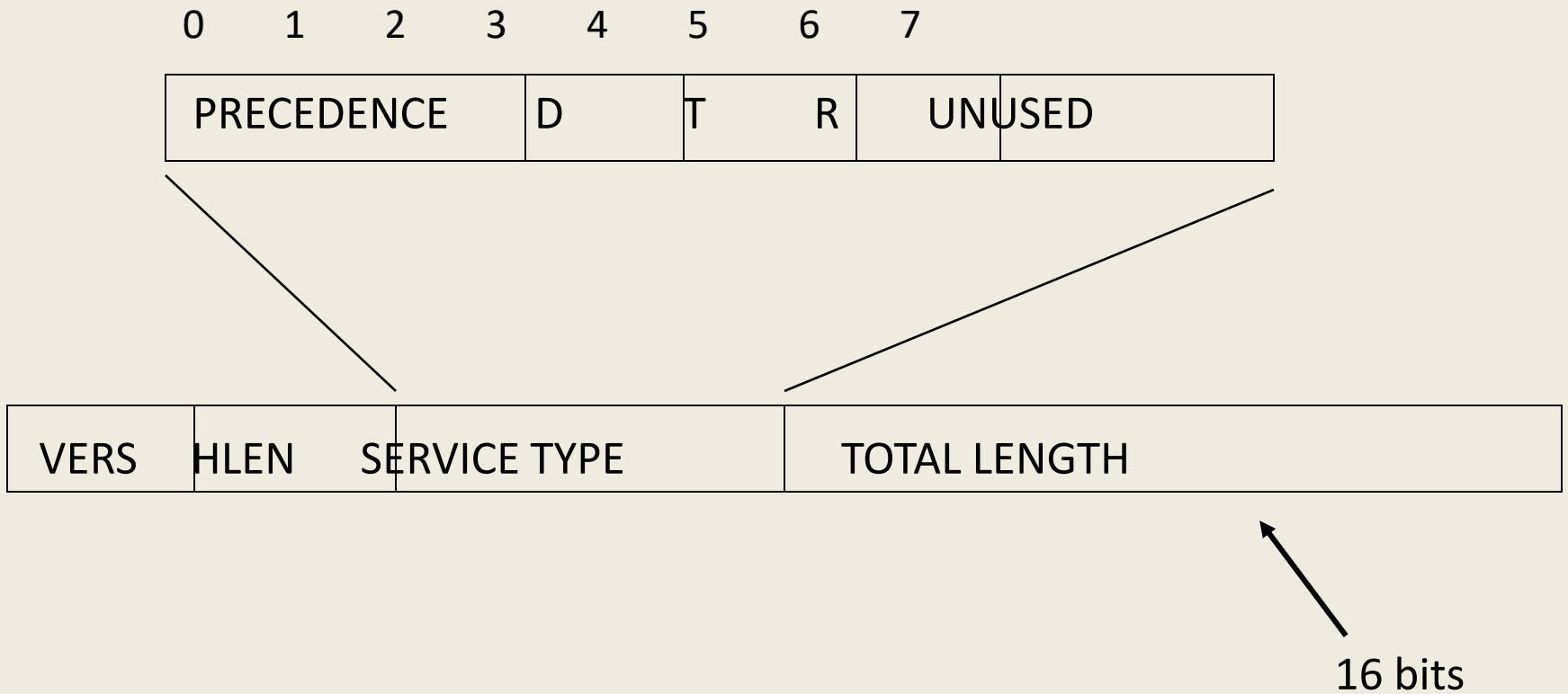


Internet

# IP Datagram Format

0	4	8	16	19	24	31
VERS	HLEN	SERVICE TYPE	TOTAL LENGTH			
IDENTIFICATION			FLAGS	FRAGMENT OFFSET		
TIME TO LIVE	PROTOCOL		HEADER CHECKSUM			
SOURCE IP ADDRESS						
DESTINATION IP ADDRESS						
IP OPTIONS (IF ANY)					PADDING	
DATA						
...						

# IP Datagram Service Type Field



# IP Datagram Service Type Field

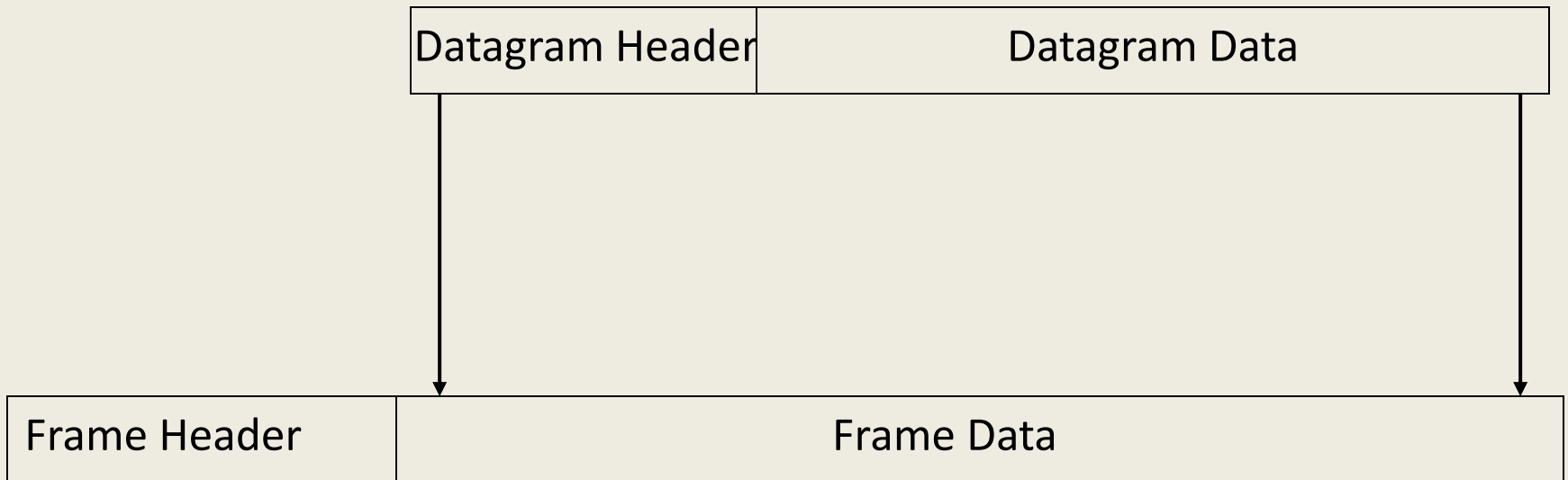
- In the 1990's the meaning of the service type field was redefined:



- For backwards compatibility:
  - When the last three codepoint bits are zero, the first three bits represent 8 classes of service (0 = least important, ..., 7 = most important)

# IP Datagram Encapsulation

- Datagrams must move from machine to machine via physical networks



- What if a datagram won't fit in a frame?



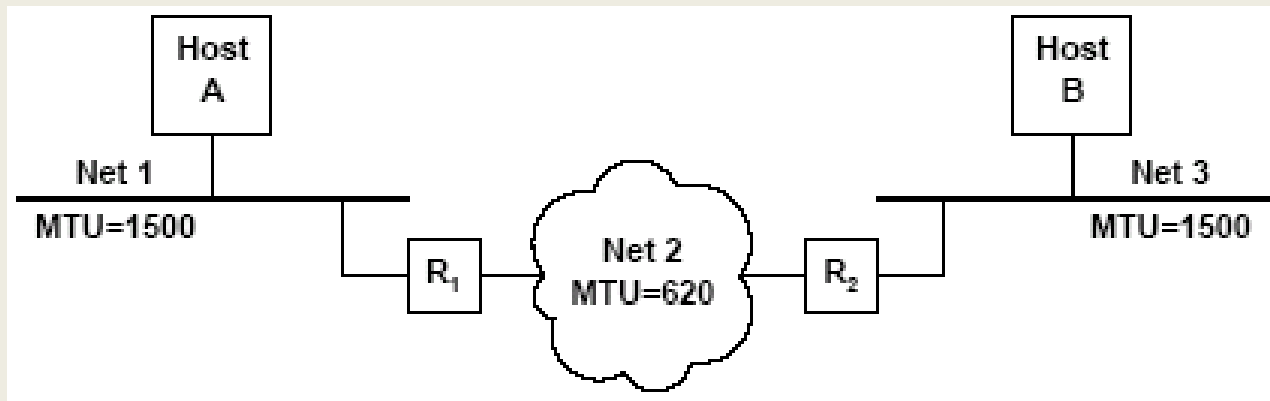
# Maximum Transfer Unit (MTU)

- Each physical networking technology limits the amount of data that can fit in a frame
  - Ethernet: 1500 octets
  - FDDI: 4470 octets
- This is called the network's MTU
- Limiting datagrams to fit in the smallest possible MTU would make travelling across networks with a larger MTU inefficient
- Allowing datagrams to be larger than a network's MTU means that datagrams will not always fit in a single frame

# Datagram Fragmentation

- Divide datagrams too large to fit in one frame into pieces called **fragments**
- Each fragment should fit into one frame
- Transport fragments over physical network
- Reassemble fragments into the complete datagram

# Datagram Fragmentation (cont)



# Datagram Fragmentation Example

Datagram Header	Data (4000 octets)
-----------------	--------------------

Fragment 1 Header	Data (1480 octets)
-------------------	--------------------

Fragment 2 Header	Data (1480 octets)
-------------------	--------------------

Fragment 3 Header	Data (1040 octets)
-------------------	--------------------

- Fragment headers duplicate datagram's header (except for FLAGS field)

# Datagram Format

0	4	8	16	19	24	31
<b>VERS</b>	<b>HLEN</b>	<b>SERVICE TYPE</b>	<b>TOTAL LENGTH</b>			
<b>IDENTIFICATION</b>			<b>FLAGS</b>	<b>FRAGMENT OFFSET</b>		
<b>TIME TO LIVE</b>		<b>PROTOCOL</b>	<b>HEADER CHECKSUM</b>			
<b>SOURCE IP ADDRESS</b>						
<b>DESTINATION IP ADDRESS</b>						
<b>IP OPTIONS (IF ANY)</b>					<b>PADDING</b>	
<b>DATA</b>						
...						

Flags - DNI, FRAG, MH

# Reassembly of Fragments

- Should fragments be reassembled after crossing one network or should we wait until the fragments arrive at their final destination to reassemble them?
- Advantages?
- Disadvantages?

# Internet Datagram Options

- Used for network testing and debugging
- Variable length field (depending on what options are selected)

0    1    2    3    4    5    6    7

Copy	Class	Option number
------	-------	---------------

- **Classes:**

- 0: Datagram or network control      1: Reserved for future use
- 2: Debugging and measurement      3: Reserved for future use

# Internet Datagram Options (cont)

0    1    2    3    4    5    6    7

Copy	Class	Option number
------	-------	---------------

- **Class 0 option numbers:**

- 0: End of options list
- 1: No operation
- 2: Security and handling restrictions
- 3: Loose source routing
- 7: Record route
- 8: Stream identifier
- 9: Strict source routing

- **Class 2 option numbers:**

- 4: Internet timestamp



# The Record Route Option

- Source creates an empty list of IP addresses in the header
- Set Record route option
- Each router that handles the datagram appends its IP address to the list
- Destination machine can extract and process the route information

# Source Route Option

- Sender dictates a path through the internet over which the datagram **must** travel
- Sender lists IP addresses (in order) of the route the datagram should take
- Sender sets the source route option

# Source Route Option (cont)

- Strict source routing - the path between two successive addresses in the list must consist of a single physical network
- Loose source routing
  - The datagram must follow the sequence of IP addresses in the list
  - Allows multiple network hops between successive addresses on the list

# The Timestamp Option

- Sender creates an empty list in the header
- Sender enables the timestamp option
- Each router that handles the datagram appends to the list its:
  - IP address
  - Local timestamp (in Universal Time)

# Processing Options During Fragmentation

0 1 2 3 4 5 6 7

Copy	Class	Option number
------	-------	---------------

- Class 0 option numbers:

- 0: End of options list
- 1: No operation
- 2: Security and handling restrictions
- 3: Loose source routing
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- Class 2 option numbers:

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# Processing Options During Fragmentation (cont)

- When fragmenting a datagram a router:
  - Replicates some IP options in all fragments
    - Example?
  - Replicates some IP options in only one fragment
    - Example?

# IP – Security Issues

- Destination IP address is used to route a datagram to its final destination
- Source IP address identifies the sender so that the receiver knows where to send a reply
- **IP spoofing** – sender of a datagram inserts the address of another machine (or a nonexistent machine) in the source address field
  - Prevent the receiver from determining the host from which an attack datagram originated
  - Make the message appear to have originated from a trusted machine
  - Want reply sent to a another (victim) host

# Teardrop

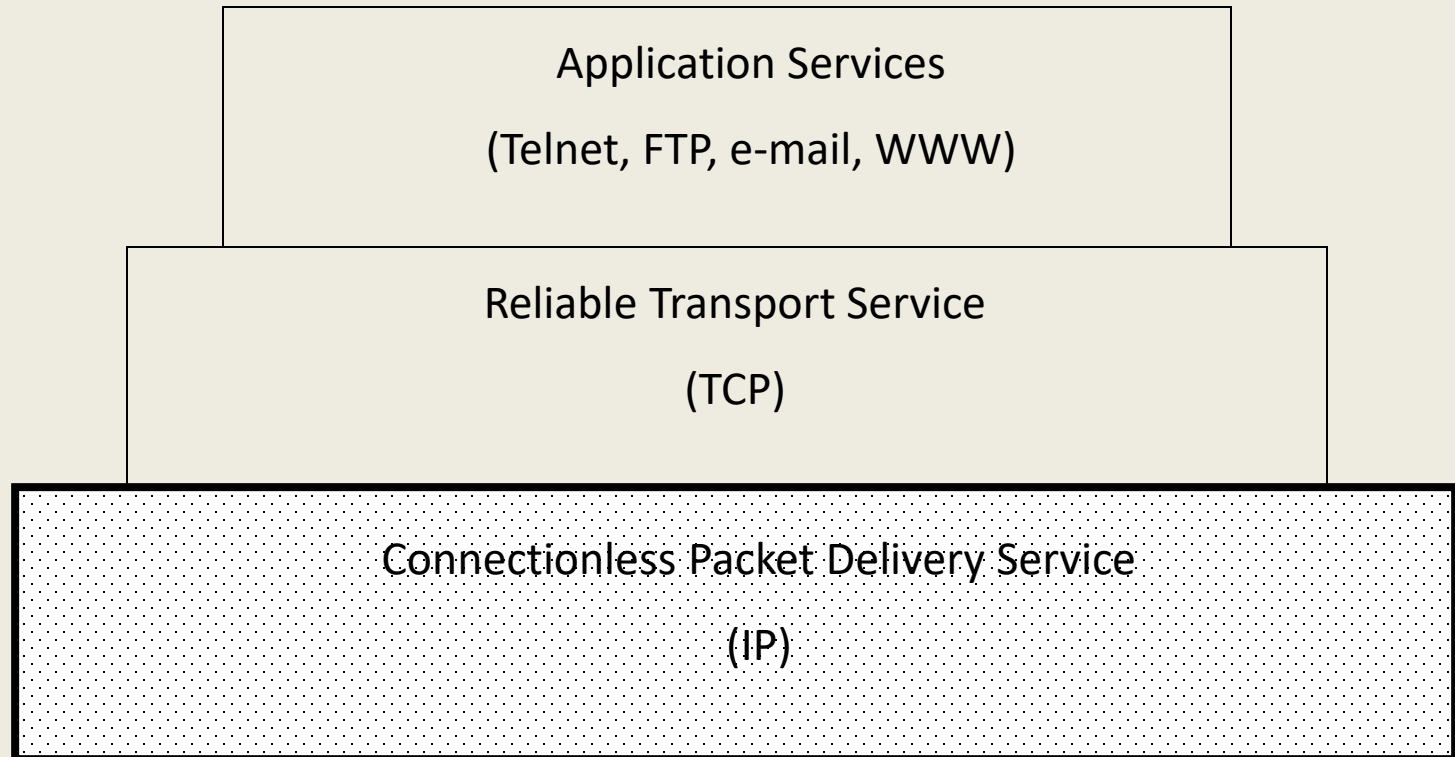
- Tool enabled attackers to crash vulnerable remote systems by sending a certain type of fragmented IP datagram
  - Normal datagram fragments do not overlap
  - Teardrop created fragments that did overlap
  - Some implementations of the TCP/IP IP fragmentation re-assembly code do not properly handle overlapping IP fragments
    - Windows and some Linux kernels
  - Caused system to crash
  - Fixed by software patches



# Summary

- The Internet Protocol (IP):
  - Provides a packet delivery service which is:
    - Unreliable
    - Best-effort
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  - Defines the basic unit of data transfer
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# Summary (cont)





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