



0/1 Knapsack Problem-

In 0/1 Knapsack Problem,

- As the name suggests, items are indivisible here.
- We can not take the fraction of any item.
- We have to either take an item completely or leave it completely.
- It is solved using dynamic programming approach.

0/1 Knapsack Problem Using Dynamic Programming-

Consider-

- Knapsack weight capacity = w
- Number of items each having some weight and value = n

0/1 knapsack problem is solved using dynamic programming in the following steps-

Step-01:

- Draw a table say 'T' with $(n+1)$ number of rows and $(w+1)$ number of columns.
- Fill all the boxes of 0^{th} row and 0^{th} column with zeroes as shown-

	0	1	2	3	W
0	0	0	0	0	0
1	0					
2	0					
.....					
n	0					

T-Table

stack and put all of its adjacent nodes which are not visited in the stack.

Step-02:

Start filling the table row wise top to bottom from left to right.



Use the following formula-

$$T(i, j) = \max \{ T(i-1, j), \text{value}_i + T(i-1, j - \text{weight}_i) \}$$

Here, $T(i, j)$ = maximum value of the selected items if we can take items 1 to i and have weight restrictions of j .

- This step leads to completely filling the table.
- Then, value of the last box represents the maximum possible value that can be put into the knapsack.

Step-03:

To identify the items that must be put into the knapsack to obtain that maximum profit,

- Consider the last column of the table.
- Start scanning the entries from bottom to top.
- On encountering an entry whose value is not same as the value stored in the entry immediately above it, mark the row label of that entry.

- After all the entries are scanned, the marked labels represent the items that must be put into the knapsack.

Time Complexity-

- Each entry of the table requires constant time $\theta(1)$ for its computation.
- It takes $\theta(nw)$ time to fill $(n+1)(w+1)$ table entries.
- It takes $\theta(n)$ time for tracing the solution since tracing process traces the n rows.
- Thus, overall $\theta(nw)$ time is taken to solve 0/1 knapsack problem using dynamic programming.

PRACTICE PROBLEM BASED ON 0/1 KNAPSACK PROBLEM-

Problem-

For the given set of items and knapsack capacity = 5 kg, find the optimal solution for the 0/1 knapsack problem making use of dynamic programming approach. ☹

Item	Weight	Value
1	2	3
2	3	4
3	4	5
4	5	6

Find the optimal solution for the 0/1 knapsack problem making use of dynamic programming approach. Consider-

$n = 4$

$w = 5 \text{ kg}$

$(w_1, w_2, w_3, w_4) = (2, 3, 4, 5)$

$(b_1, b_2, b_3, b_4) = (3, 4, 5, 6)$

OR

A thief enters a house for robbing it. He can carry a maximal weight of 5 kg into his bag. There are 4 items in the house with the following weights and values. What items should thief take if he either takes the item completely or leaves it completely?

Item	Weight (kg)	Value (\$)
Mirror	2	3
Silver nugget	3	4
Painting	4	5
Vase	5	6

Solution-

Given-

- Knapsack capacity (w) = 5 kg
- Number of items (n) = 4

Step-01:

- Draw a table say 'T' with $(n+1) = 4 + 1 = 5$ number of rows and $(w+1) = 5 + 1 = 6$ number of columns.
- Fill all the boxes of 0^{th} row and 0^{th} column with 0.

	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0					
2	0					
3	0					
4	0					

T-Table

Step-02:

Start filling the table row wise top to bottom from left to right using the formula-

$$T(i, j) = \max \{ T(i-1, j), \text{value}_i + T(i-1, j - \text{weight}_i) \}$$

Finding T(1,1)-

We have,

- $i = 1$
- $j = 1$
- $(\text{value})_i = (\text{value})_1 = 3$
- $(\text{weight})_i = (\text{weight})_1 = 2$

Substituting the values, we get-

$$T(1,1) = \max \{ T(1-1, 1), 3 + T(1-1, 1-2) \}$$

$$T(1,1) = \max \{ T(0,1), 3 + T(0,-1) \}$$

$$T(1,1) = T(0,1) \quad \{ \text{Ignore } T(0,-1) \}$$

$$T(1,1) = 0$$

Finding T(1,2)-

We have,

- $i = 1$
- $j = 2$
- $(\text{value})_i = (\text{value})_1 = 3$
- $(\text{weight})_i = (\text{weight})_1 = 2$

Substituting the values, we get-

$$T(1,2) = \max \{ T(1-1, 2), 3 + T(1-1, 2-2) \}$$

$$T(1,2) = \max \{ T(0,2), 3 + T(0,0) \}$$

$$T(1,2) = \max \{ 0, 3+0 \}$$

$$T(1,2) = 3$$

Finding $T(1,3)$ -

We have,

- $i = 1$
- $j = 3$
- $(\text{value})_i = (\text{value})_1 = 3$
- $(\text{weight})_i = (\text{weight})_1 = 2$

Substituting the values, we get-

$$T(1,3) = \max \{ T(1-1, 3), 3 + T(1-1, 3-2) \}$$

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$$T(1,3) = \max \{ T(0,3), 3 + T(0,1) \}$$

$$T(1,3) = \max \{ 0, 3+0 \}$$

$$T(1,3) = 3$$

Finding $T(1,4)$ -

We have,

- $i = 1$

- $j = 4$
- $(value)_i = (value)_1 = 3$
- $(weight)_i = (weight)_1 = 2$

Substituting the values, we get-

$$T(1,4) = \max \{ T(1-1, 4), 3 + T(1-1, 4-2) \}$$

$$T(1,4) = \max \{ T(0,4), 3 + T(0,2) \}$$

$$T(1,4) = \max \{ 0, 3+0 \}$$

$$T(1,4) = 3$$



Finding $T(1,5)$ -

We have,

- $i = 1$
- $j = 5$
- $(value)_i = (value)_1 = 3$
- $(weight)_i = (weight)_1 = 2$

Substituting the values, we get-

$$T(1,5) = \max \{ T(1-1, 5), 3 + T(1-1, 5-2) \}$$

$$T(1,5) = \max \{ T(0,5), 3 + T(0,3) \}$$

$$T(1,5) = \max \{ 0, 3+0 \}$$

$$T(1,5) = 3$$

Finding $T(2,1)$ -

We have,

- $i = 2$
- $j = 1$
- $(value)_i = (value)_2 = 4$

- $(\text{weight})_i = (\text{weight})_2 = 3$

Substituting the values, we get-

$$T(2,1) = \max \{ T(2-1, 1), 4 + T(2-1, 1-3) \}$$

$$T(2,1) = \max \{ T(1,1), 4 + T(1,-2) \}$$

$$T(2,1) = T(1,1) \quad \{ \text{Ignore } T(1,-2) \}$$

$$T(2,1) = 0$$

Finding $T(2,2)$ -

We have,

- $i = 2$
- $j = 2$
- $(\text{value})_i = (\text{value})_2 = 4$
- $(\text{weight})_i = (\text{weight})_2 = 3$

Substituting the values, we get-

$$T(2,2) = \max \{ T(2-1, 2), 4 + T(2-1, 2-3) \}$$

$$T(2,2) = \max \{ T(1,2), 4 + T(1,-1) \}$$

$$T(2,2) = T(1,2) \quad \{ \text{Ignore } T(1,-1) \}$$

$$T(2,2) = 3$$

Finding $T(2,3)$ -

We have,

- $i = 2$
- $j = 3$
- $(\text{value})_i = (\text{value})_2 = 4$
- $(\text{weight})_i = (\text{weight})_2 = 3$

Substituting the values, we get-

$$T(2,3) = \max \{ T(2-1, 3), 4 + T(2-1, 3-3) \}$$

$$T(2,3) = \max \{ T(1,3), 4 + T(1,0) \}$$

$$T(2,3) = \max \{ 3, 4+0 \}$$

$$T(2,3) = 4$$

Finding T(2,4)-

We have,

- $i = 2$
- $j = 4$
- $(\text{value})_i = (\text{value})_2 = 4$
- $(\text{weight})_i = (\text{weight})_2 = 3$

Substituting the values, we get-

$$T(2,4) = \max \{ T(2-1, 4), 4 + T(2-1, 4-3) \}$$

$$T(2,4) = \max \{ T(1,4), 4 + T(1,1) \}$$

$$T(2,4) = \max \{ 3, 4+0 \}$$

$$T(2,4) = 4$$

Finding T(2,5)-

We have,

- $i = 2$
- $j = 5$
- $(\text{value})_i = (\text{value})_2 = 4$
- $(\text{weight})_i = (\text{weight})_2 = 3$

Substituting the values, we get-

$$T(2,5) = \max \{ T(2-1, 5), 4 + T(2-1, 5-3) \}$$

$$T(2,5) = \max \{ T(1,5), 4 + T(1,2) \}$$

$$T(2,5) = \max \{ 3, 4+3 \}$$

$$T(2,5) = 7$$

Similarly, compute all the entries.

After all the entries are computed and filled in the table, we get the following table-

	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0	0	3	3	3	3
2	0	0	3	4	4	7
3	0	0	3	4	5	7
4	0	0	3	4	5	7

T-Table

- The last entry represents the maximum possible value that can be put into the knapsack.
- So, maximum possible value that can be put into the knapsack = 7.

Identifying Items To Be Put Into Knapsack-

Following Step-04,

- We mark the rows labelled "1" and "2".
- Thus, items that must be put into the knapsack to obtain the maximum value 7 are-

Item-1 and Item 2