



SNS COLLEGE OF ENGINEERING

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An Autonomous Institution

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University, Chennai

DEPARTMENT OF INFORMATION TECHNOLOGY

**Course Code and Name : 19IT602– CRYPTOGRAPHY AND CYBER
SECURITY**

III YEAR / VI SEMESTER

**Unit 4: CYBER SECURITY VULNERALIBILITES AND MESSAGE
AUTHENTICATION**

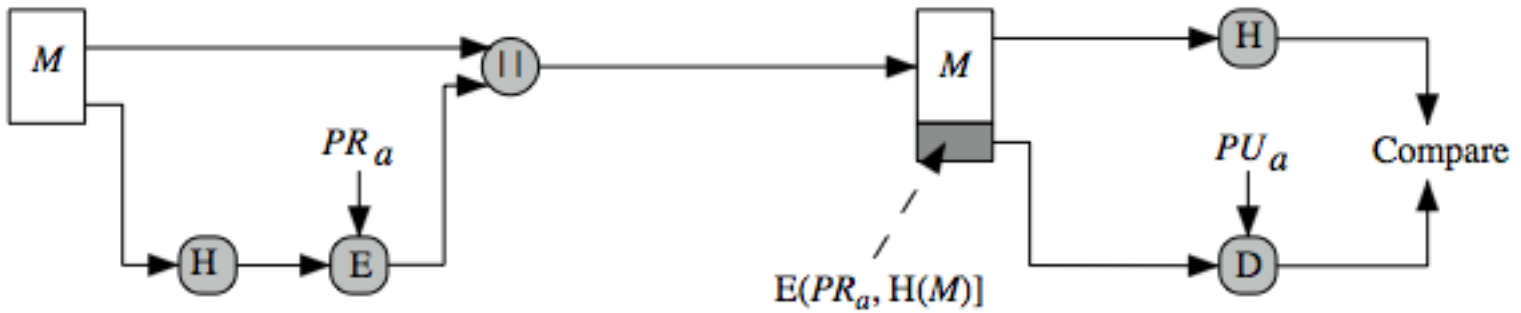
Topic : DSS



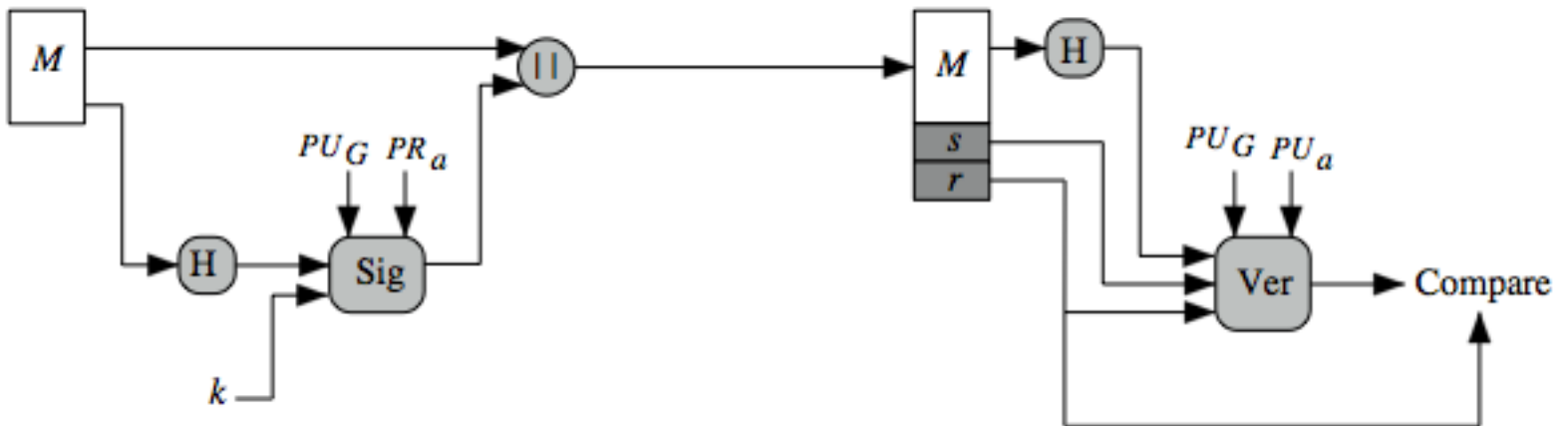
Digital Signature Standard (DSS)

- US Govt approved signature scheme
- designed by NIST & NSA in early 90's
- published as FIPS-186 in 1991
- revised in 1993, 1996 & then 2000
- uses the SHA hash algorithm
- DSS is the standard, DSA is the algorithm
- FIPS 186-2 (2000) includes alternative RSA & elliptic curve signature variants
- DSA is digital signature only unlike RSA
- is a public-key technique

DSS vs RSA Signatures



(a) RSA Approach



(b) DSS Approach



Digital Signature Algorithm (DSA)

- creates a 320 bit signature
- with 512-1024 bit security
- smaller and faster than RSA
- a digital signature scheme only
- security depends on difficulty of computing discrete logarithms
- variant of ElGamal & Schnorr schemes



DSA Key Generation

- have shared global public key values (p, q, g) :
 - choose 160-bit prime number q
 - 160 bit prime divisor of $(p-1)$ $2^{159} < q < 2^{160}$
 - choose a large prime p with $2^{L-1} < p < 2^L$
 - where $L = 512$ to 1024 bits and is a multiple of 64
 - choose $g = h^{(p-1)/q}$
 - where $1 < h < p-1$ and $h^{(p-1)/q} \bmod p > 1$
 - Must be greater than 1
- users choose private & compute public key:
 - choose random private key: $x < q$
 - compute public key: $y = g^x \bmod p$



DSA Signature Creation

➤ to **sign** a message M the sender:

- generates a random signature key k , $k < q$
- nb. k must be random, be destroyed after use, and never be reused

➤ then computes signature pair:

$$r = (g^k \bmod p) \bmod q$$
$$s = [k^{-1} (H(M) + xr)] \bmod q$$

➤ sends signature (r, s) with message M



DSA Signature Verification

- having received M & signature (r, s)
- to **verify** a signature, recipient computes:

$$w = s^{-1} \bmod q$$

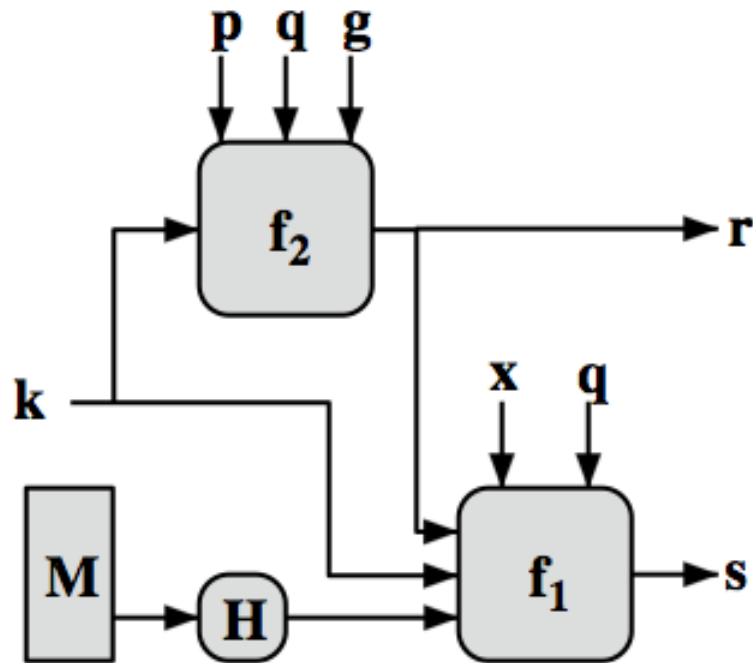
$$u1 = [H(M)w] \bmod q$$

$$u2 = (rw) \bmod q$$

$$v = [(g^{u1} y^{u2}) \bmod p] \bmod q$$

- if $v=r$ then signature is verified

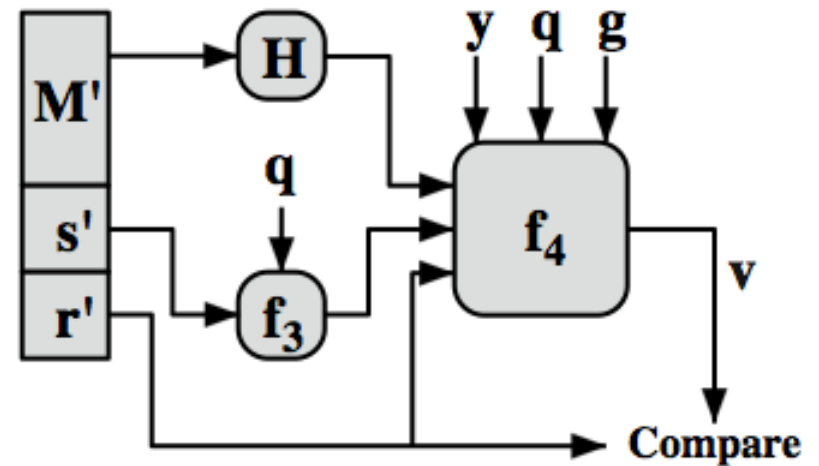
DSS Overview



$$s = f_1(H(M), k, x, r, q) = (k^{-1} (H(M) + xr)) \bmod q$$

$$r = f_2(k, p, q, g) = (g^k \bmod p) \bmod q$$

(a) Signing



$$w = f_3(s', q) = (s')^{-1} \bmod q$$

$$v = f_4(y, q, g, H(M'), w, r')$$

$$= ((g^{(H(M')w) \bmod q} y^{r'w \bmod q}) \bmod p) \bmod q$$

(b) Verifying



Thank You