



SNS COLLEGE OF ENGINEERING

Coimbatore-107



COURSE NAME: ANALYSIS OF ALGORITHM

II YEAR/ IV SEMESTER

UNIT – III

DYNAMIC PROGRAMMING

Topic

Optimal Binary Search Tree



Dynamic Programming Optimal Binary Search Tree (OBST)

→ It is an efficient searching algorithm, that works on sorted arrays.

→ It minimizes total search cost considering the frequency of searches [number of times key is searched] for each element.

Example: (A sorted Array)

Key	10	20	30	40	50
Freq	1	5	10	3	2

Step 1:

pick the element with maximum frequency & make it root.

30

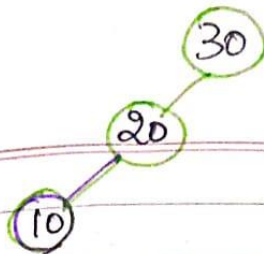
Step 2:

→ Build Left & Right Subtrees with key elements 10 and 20 (left) and 40, 50 (as right).

Step 3: Construct Left subtree with key 10 and 20.

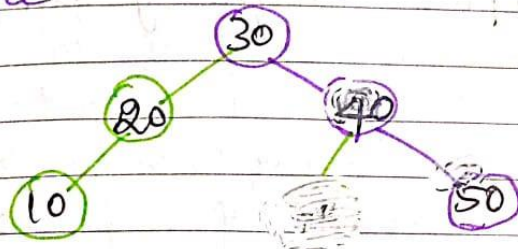
By Rule of 'BST' - 10 and 20 are smaller than root. So construct it as Left subtree.

Also frequency of 20 is 5 ^{close} to root. So take it as ^{Immediate} left child.



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Step 4 : Construct Right subtree
with key 40 and 50.
By Rule of BST - 40 and 50 are
greater. So construct it as Right
subtree



Construct OBST Using Dynamic
programming. (for same problem)

We use Dynamic programming to
minimize the search cost and
construct tree.

Step 1 : Sort the keys in Ascending order

Step 2 : calculate minimum cost of
Searching for all possible subtrees.

Step 3 : Build optimal tree.

Define the cost :

(i) $C[i, j]$ \rightarrow Minimum search cost from
index i to j .

(ii) $W[i, j]$ \rightarrow Sum of frequencies from i to j .



Recurrence Relation:

Cost of subtree with root 'k' can be given as

$$C[i, j] = \min_{i \leq k \leq j} [C(i, k-1) + C(k+1, j) + W(i, j)]$$

Goal:

To find 'k' that minimizes the cost for each subtree

Step 4: Compute Frequency Sum Table $W[i, j]$

- 1 {
 - (i) $W(1, 1) = \text{Frequency of } 10 = 1$
 - (ii) $W(1, 2) = \text{Frequency of } 10 \text{ \& } 20 = 1 + 5 = 6$
 - (iii) $W(1, 3) = \text{Frequency of } 10, 20, 30 = 1 + 5 + 10 = 16$
 - (iv) $W(1, 4) = \text{Frequency of } 10, 20, 30, 40 = 1 + 5 + 10 + 3 = 19$
 - (v) $W(1, 5) = \text{Freq. of } 10, 20, 30, 40, 50 = 1 + 5 + 10 + 3 + 2 = 21$
- 2 {
 - (vi) $W(2, 2) = \text{Freq. of } 20 = 5$
 - (vii) $W(2, 3) = \text{Freq. of } 20, 30 = 5 + 10 = 15$
 - (viii) $W(2, 4) = \text{Freq. of } 20, 30, 40 = 5 + 10 + 3 = 18$
 - (ix) $W(2, 5) = \text{Freq. of } 20, 30, 40, 50 = 5 + 10 + 3 + 2 = 20$
- 3 {
 - (x) $W(3, 3) = \text{Freq. of } 30 = 10$
 - (xi) $W(3, 4) = \text{Freq. of } 30, 40 = 10 + 3 = 13$
 - (xii) $W(3, 5) = \text{Freq. of } 30, 40, 50 = 10 + 3 + 2 = 15$
- 4 {
 - (xiii) $W(4, 4) = \text{Freq. of } 40 = 3$
 - (xiv) $W(4, 5) = \text{Freq. of } 40, 50 = 3 + 2 = 5$
- 5 {
 - (xv) $W(5, 5) = \text{Freq. of } 50 = 2$

Step 4: DP Table $C[i, j] \rightarrow$ Minimum Cost



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i	j	c[i,j]
10	10	1 (5)
20	20	5 (2)
30	30	10 (1)
40	40	3 (3)
50	50	2 (4)

OBSTree

30' as Root with highest frequency (10)

20 < 30; Left 2nd High frequency close to root (20)

10 < 30; left (10) → less frequency

30 < 40; Right; 2nd High frequency close to root (40)

30 < 50; Right (50) → less frequency

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